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Understanding DMA Malware

DIMVA2012 - 9th Conference on Detection of Intrusions and Malware & Vulnerability Assessment

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Malicious Software Arms Race

Log on

Welcome to Internet Banking

Log on to Personal Internet Banking
Please enter your user ID eg IB1234567890 or John123
 Continue

Remember my user ID [? Forgotten your user ID?](#)

New customers

▶ [Register for Internet Banking](#)
▶ [Activate your Secure Key](#)

Business customers

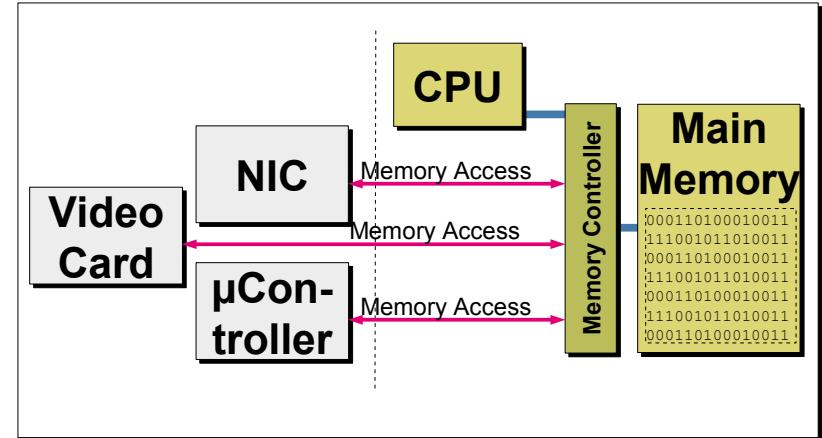
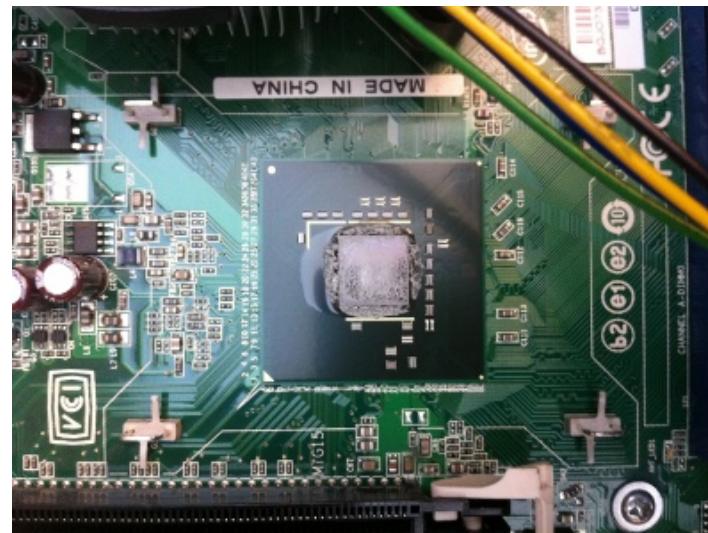
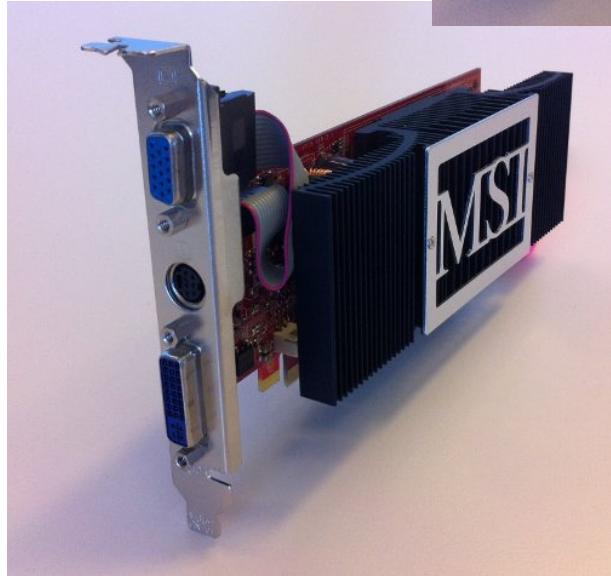
▶ [Log on to Business Internet Banking](#)
▶ [Apply](#)
▶ [Activate](#)

(www.hsbc.co.uk)

- Countermeasures
Anti-virus, firewalls, etc.
- Stealth attacks,
see *rootkit evolution*

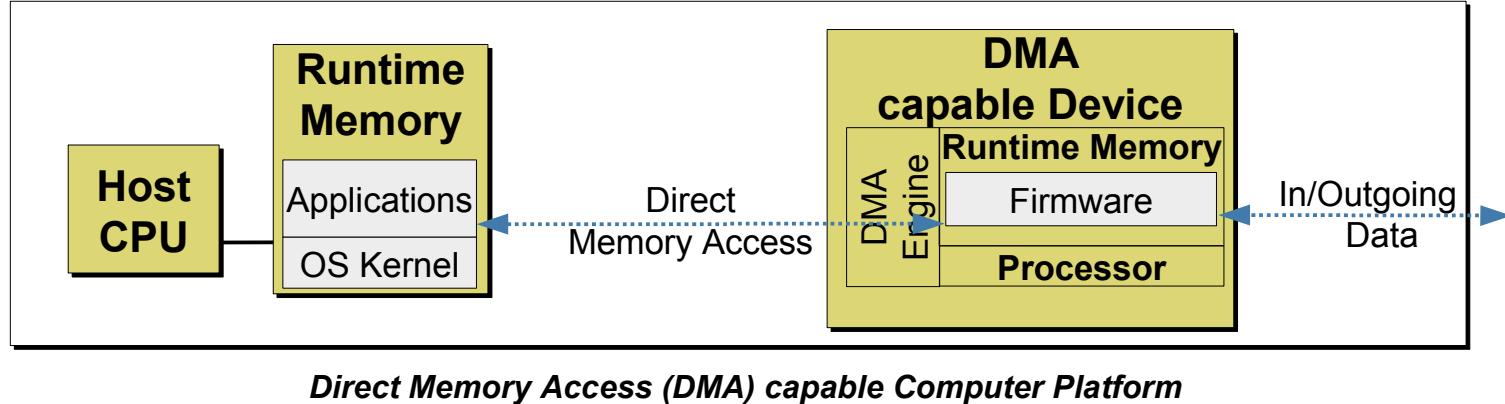


Dedicated Hardware

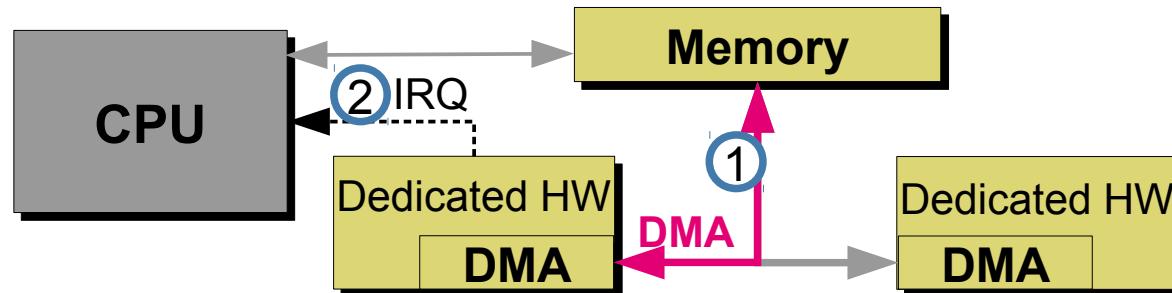


Separated Execution Environments

Common Hardware (HW) Features

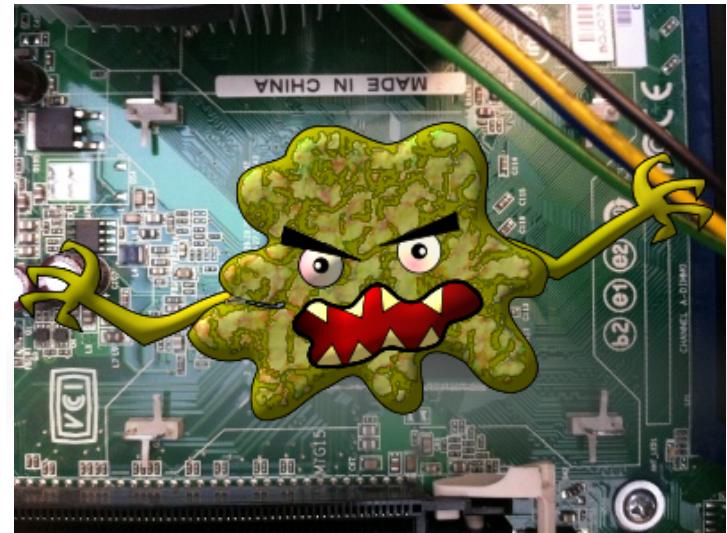


- Precondition for stealth malware → 1st party DMA:



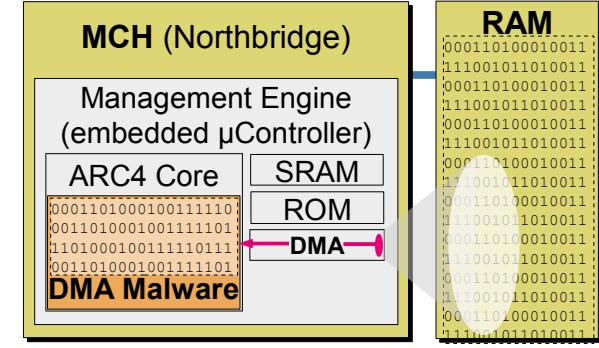
DMA Malware Definition

- **More than controlling a DMA engine**
- Malware functionality executed on dedicated HW
- No physical access
- Rootkit/stealth capabilities
- Optional:
Survival of power off mode



DMA Malware Properties

- **Three phases**
 - Search
 - Process data
 - Exfiltration/infiltration



DMA Capable Device Integrated in Intel Chipsets

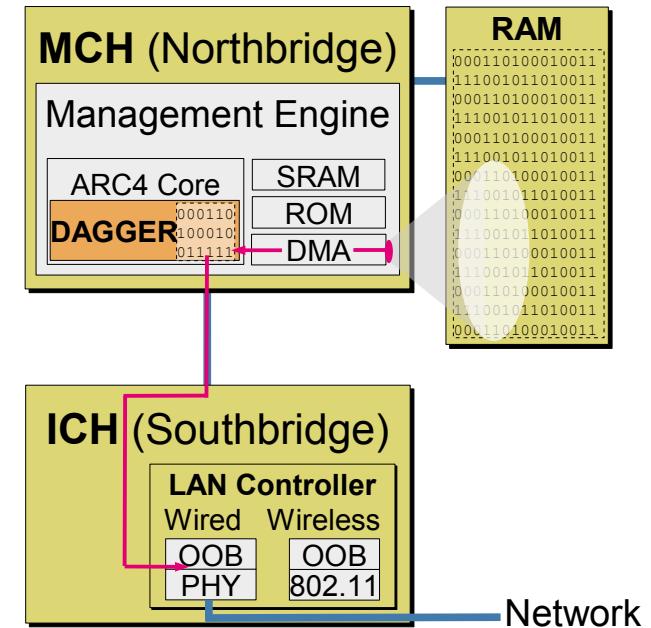
- **Core functionality**
 - Virtual/physical memory address mapping
 - Overcoming address randomization
 - Search space restriction

Comparison of DMA Attacks

	Malware Functionality	Without Physical Access	Rootkit/Stealth Capabilities	Survives Reboot/Standby/Power off
USB [Maynor '05]	no	no	no	yes
Firewire [Dornseif et al.'04/'05] & [Boileau'06]	yes	no	yes	yes
PCMCIA [Aumaitre et al.'10]	yes	no	yes	yes
ME [Tereshkin'09]	no	yes	no	(yes)
NIC [Duflot et al.'10]	yes	yes	yes	yes
NIC [Delugre'10]	yes	yes	yes	yes
Video+NIC [Triulzi'08/'10]	yes	yes	yes	yes
<i>this work</i>	yes	yes	yes	(yes)

DAGGER – Our DMA Malware Example

- *DmA based keyloGGER*
- Implements all three phases
 - Search keyboard buffer
 - Monitor keyboard buffer
 - Exfiltrate keystroke codes
- Evaluation of core functionalities
- Proof of concept for stealing short-living runtime data stealthily!
- Infiltration via security vulnerability



DAGGER Monitoring the Host's Keyboard Buffer

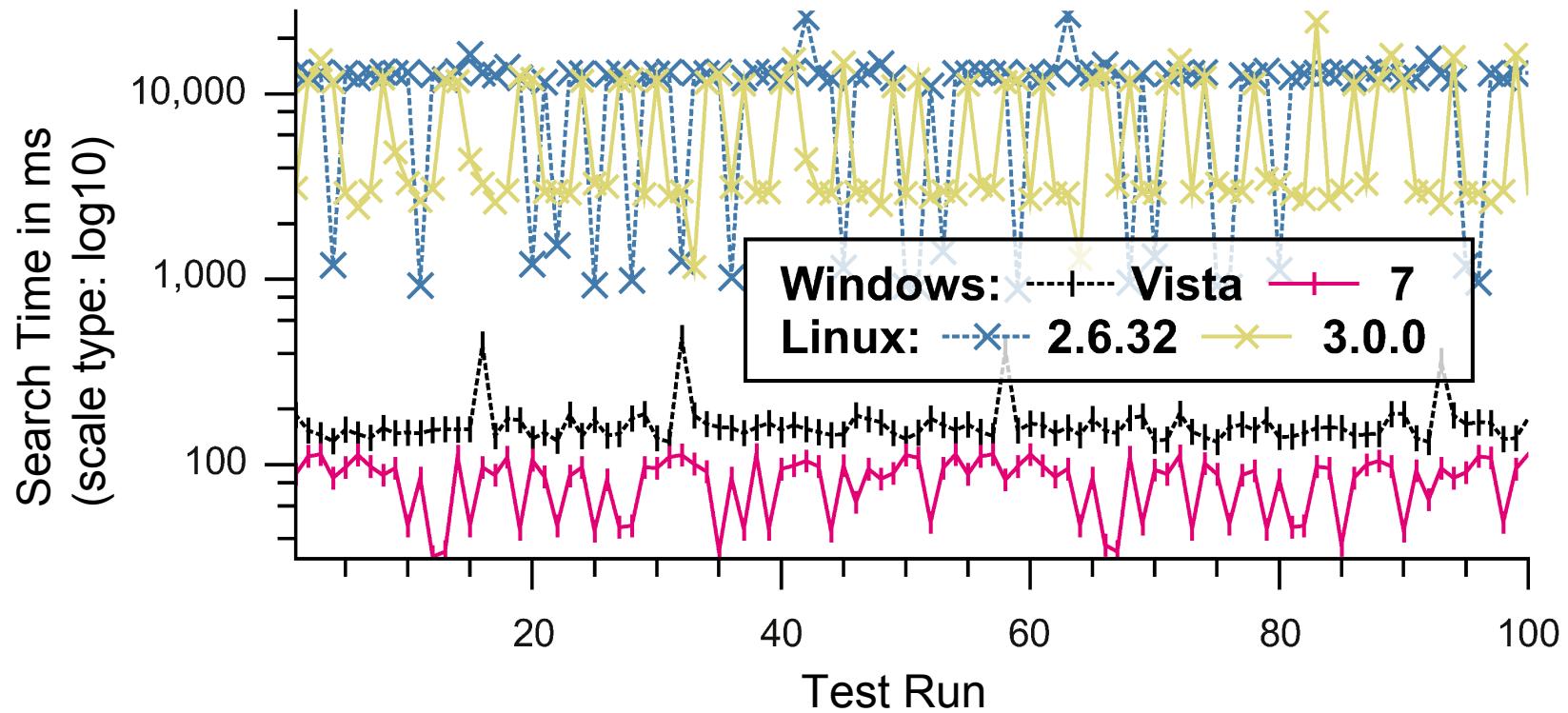
DAGGER Implementation

- Different search strategies

	Windows	Linux
Virtual/physical memory address mapping	page tables	offset
Address randomization	randomization mechanism in place	no randomization
Search space restrictions	<i>Object Manager Namespace Directory</i>	address ranges

- Platform: Intel Q35 chipset, 2GB RAM, 4-core 3GHz CPU

Evaluation – Several Operating System Kernels



Evaluation – Attacking Linux Harddisk Encryption

- *Aggressive search mode*
- Linux Unified Key Setup (LUKS)/
Device Mapper's crypt (dm-crypt) setup
- DAGGER can catch pre boot authentication
passphrase



Evaluation

- Anti-virus software, firewalls, Wireshark, Mamutu, etc.
- Several USB keyboards
- Windows swap behavior
- Performance overhead for host system
- Manageability Engine firmware condition
 - Status tools
 - Active Management Technology webserver



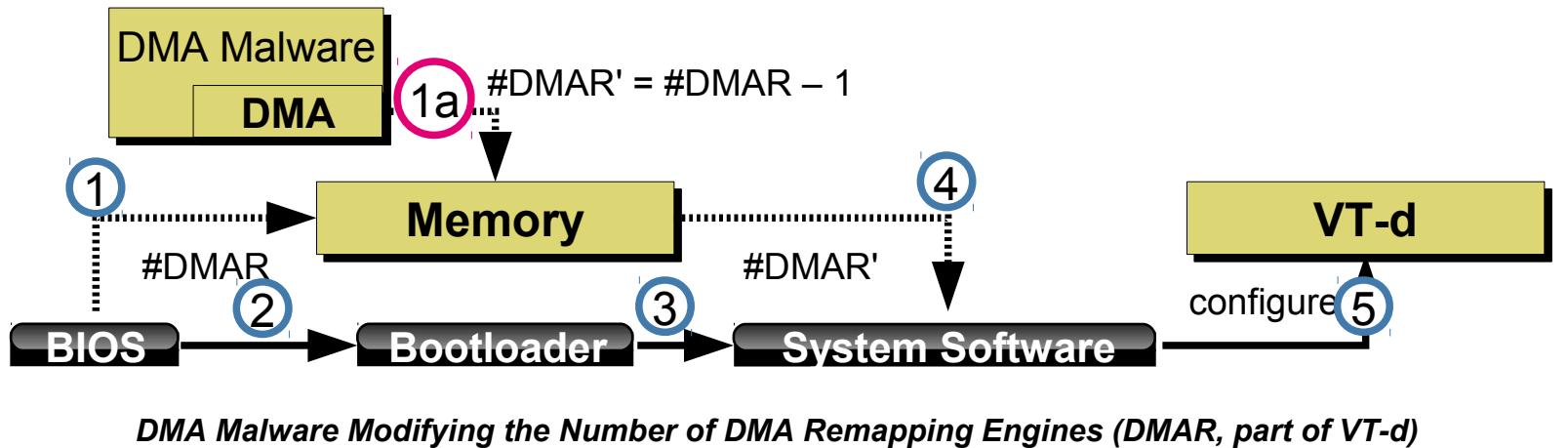
Intel® Active Management Technology
Computer: q35

System Status	
Hardware Information	
System	Off
Processor	192.1
Memory	14bfa
Disk	7/19/2
Event Log	
Remote Control	4:56 p
Power Policies	
Network Settings	
User Accounts	

Refresh

Countermeasures

- *Input/Output Memory Management Unit (I/OMMU)*
- Intel: *Virtualization Technology for Directed I/O (VT-d)*
- Issues
 - Missing (Windows) or experimental (Linux) drivers
 - *CoPilot* [Petroni et al.'04]/*DeepWatch* [Bulygin'08] or DAGGER? → policy conflict
 - Attack with DAGGER's execution environment



Conclusion

- DMA Malware definition
- Focus on stealth attacks
- Evaluation of DMA Malware core functionalities
- DMA Malware is
 - Effective
 - Efficient enough for real attacks
- Specialized countermeasures must be developed





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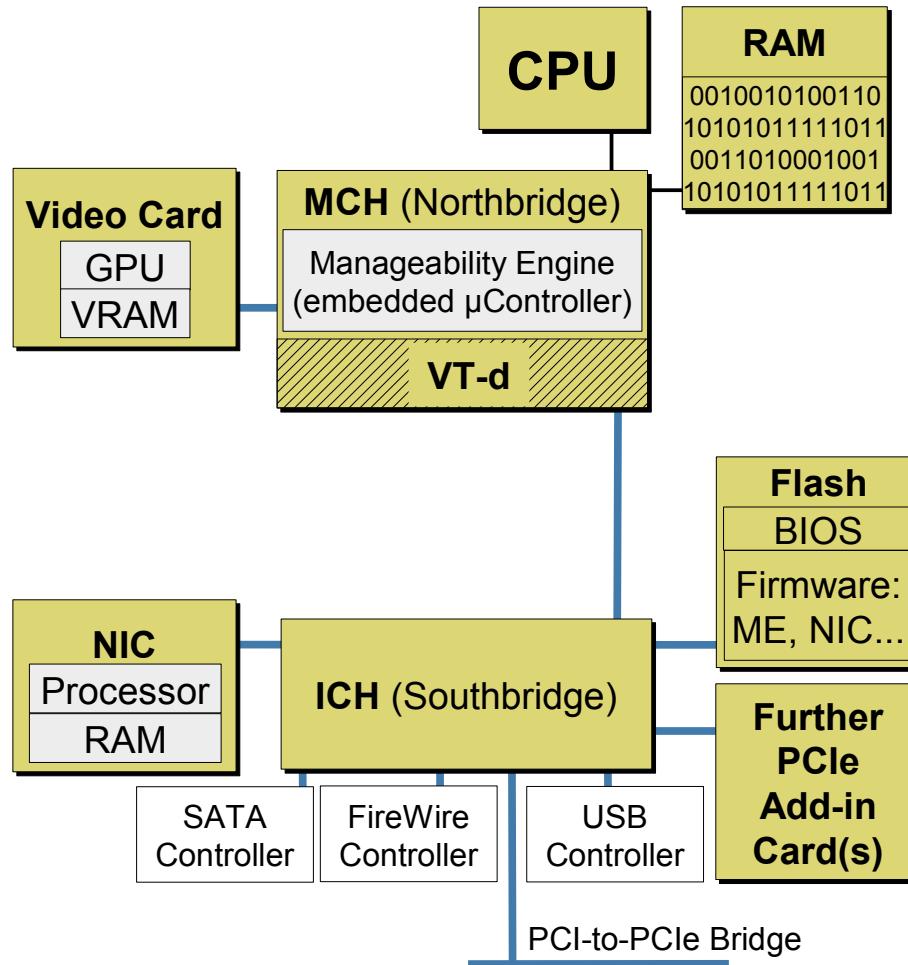
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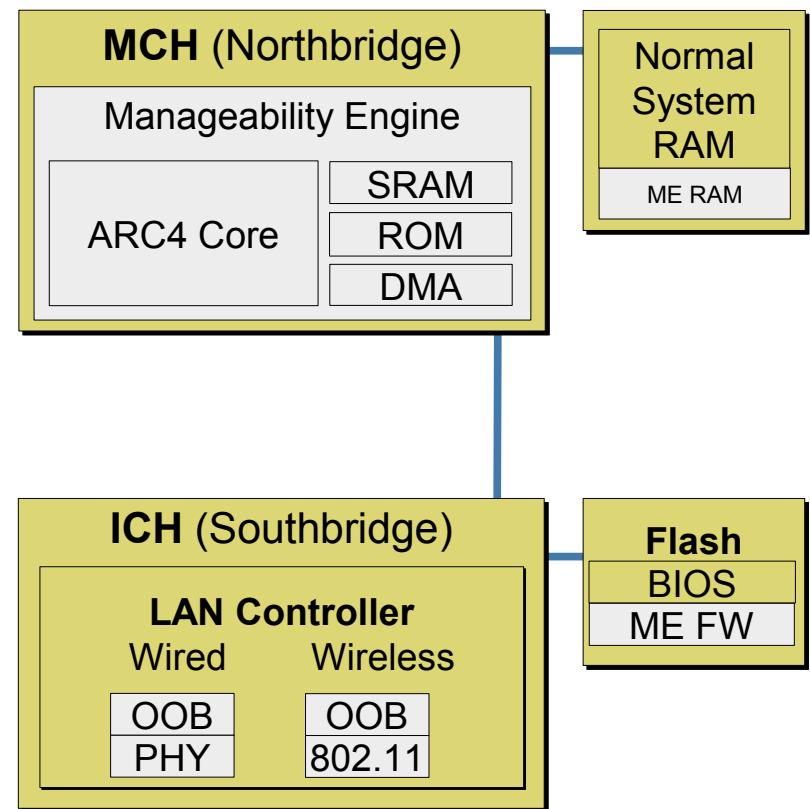
Backup

Background – x86 Platform

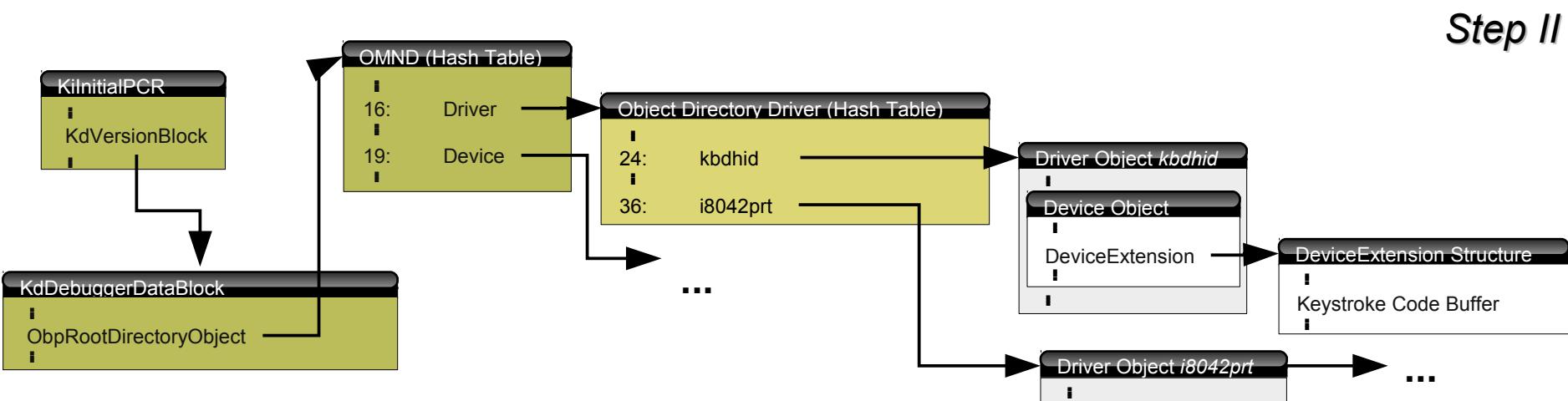
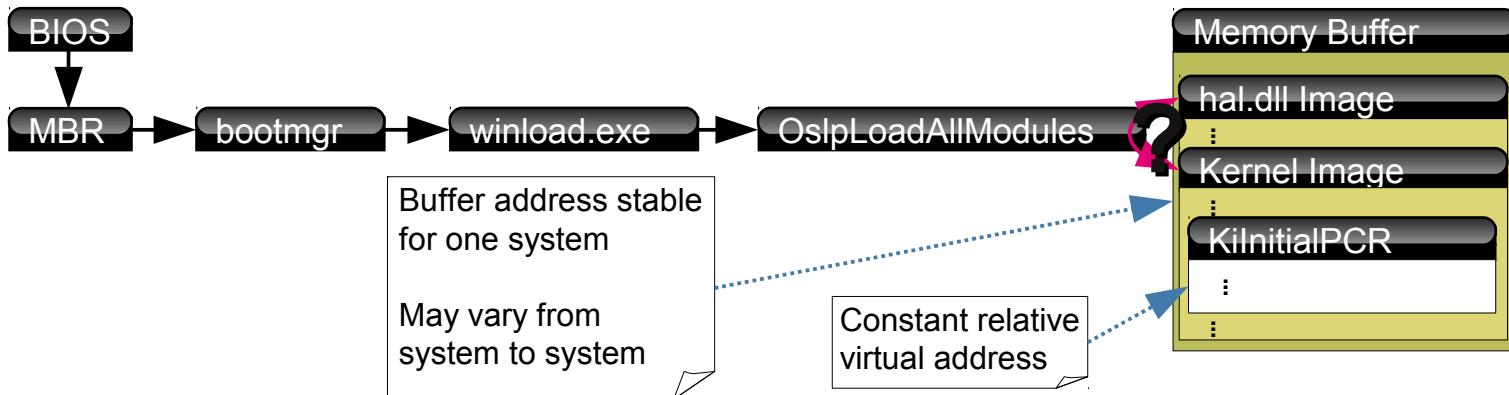


DAGGER – Example DmA based KeyloGGER Malware

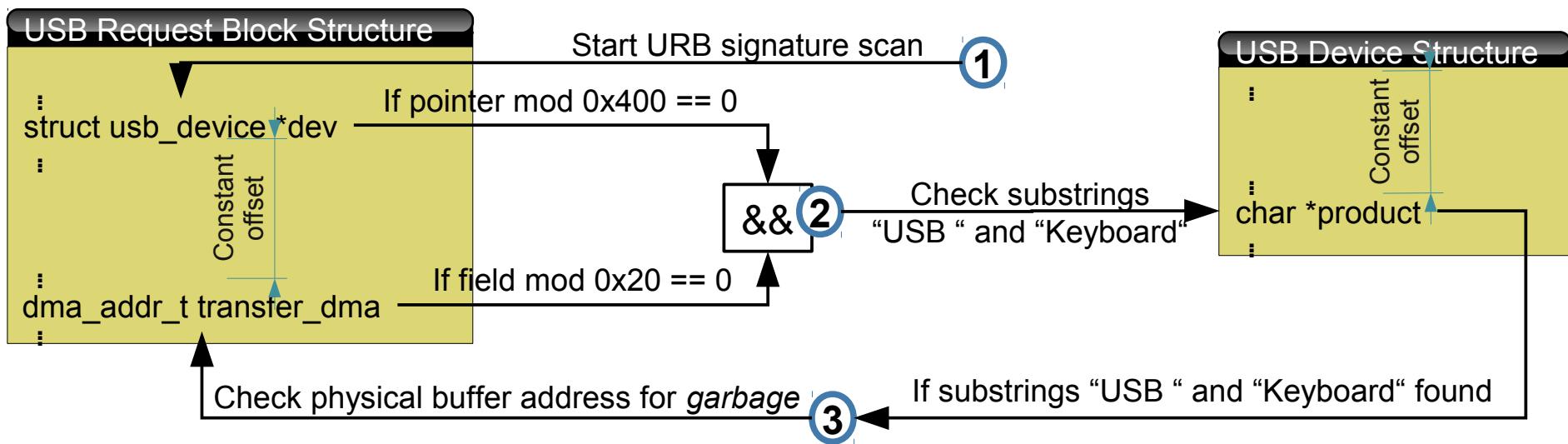
- Implementation based on Intel Manageability Engine (ME)
 - Executes firmware such as *Active Management Technology, Identity Protection Technology, Integrated Trusted Platform Module*, etc.
- Objectives
 - Find keyboard buffer
 - Permanently monitor keyboard buffer
 - Exfiltrate keystroke codes



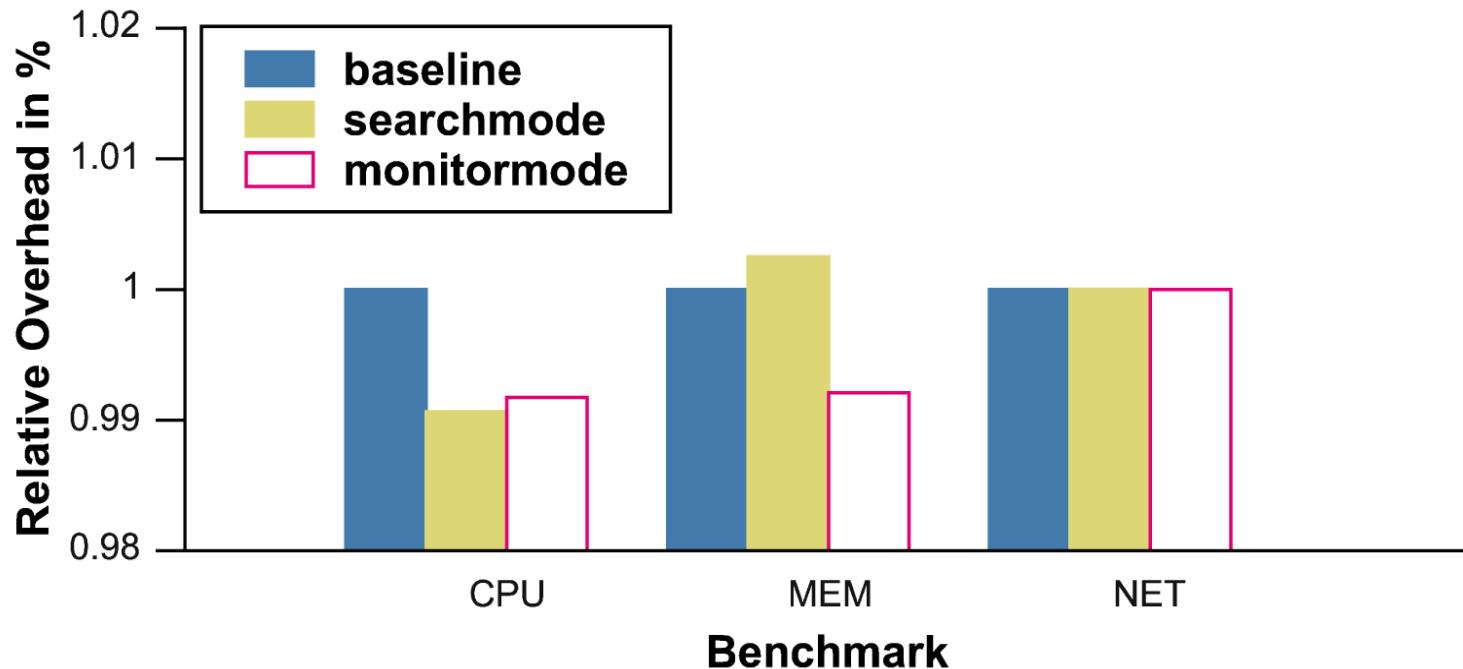
Windows Attack Details



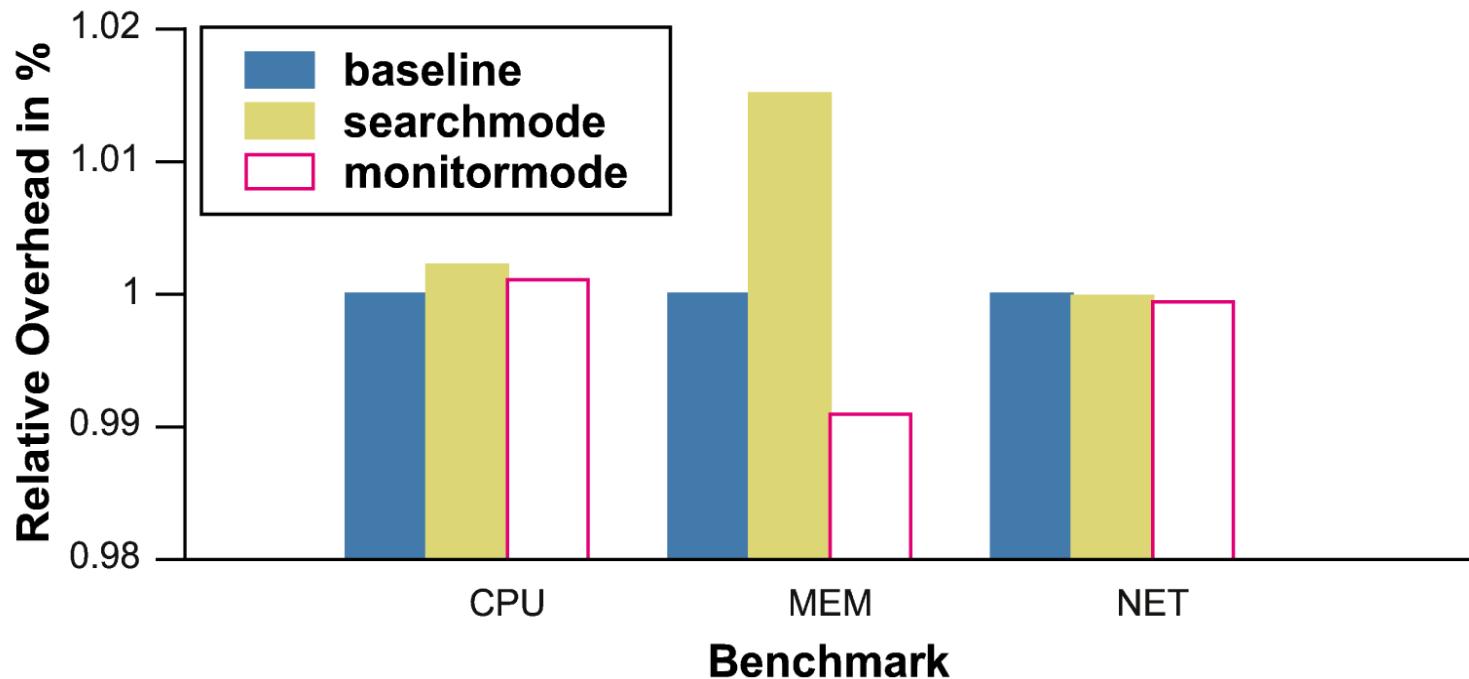
Linux Attack Details



DAGGER Evaluation – Performance Overhead, Windows Host



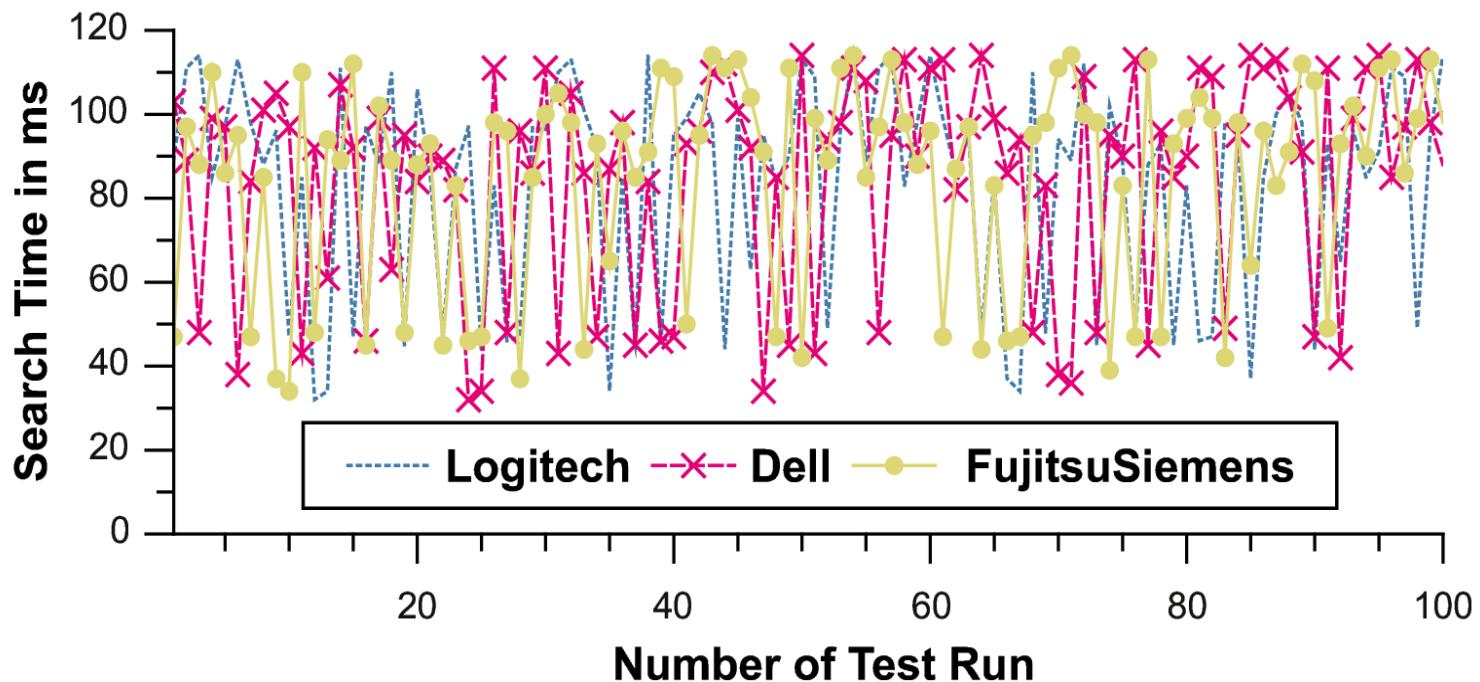
Evaluation – Performance Overhead, Linux Host



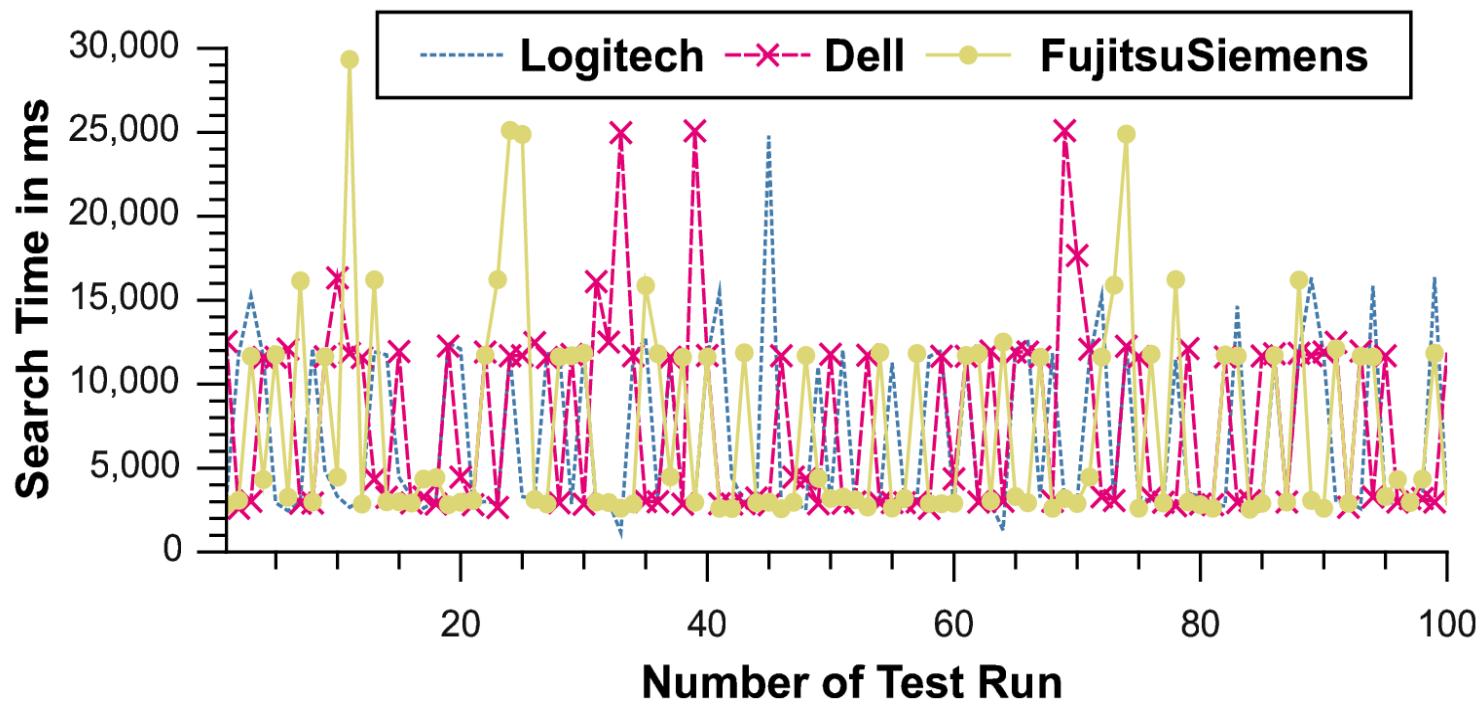
Evaluation – Effectiveness and Efficiency

- Several Operating System Kernels
 - Windows 7
 - Windows Vista
 - Linux 3.0.0
 - Linux 2.6.32
- Several Keyboards
 - Logitech
 - Dell
 - FujitsuSiemens
- Swap file behavior
 - Windows 7

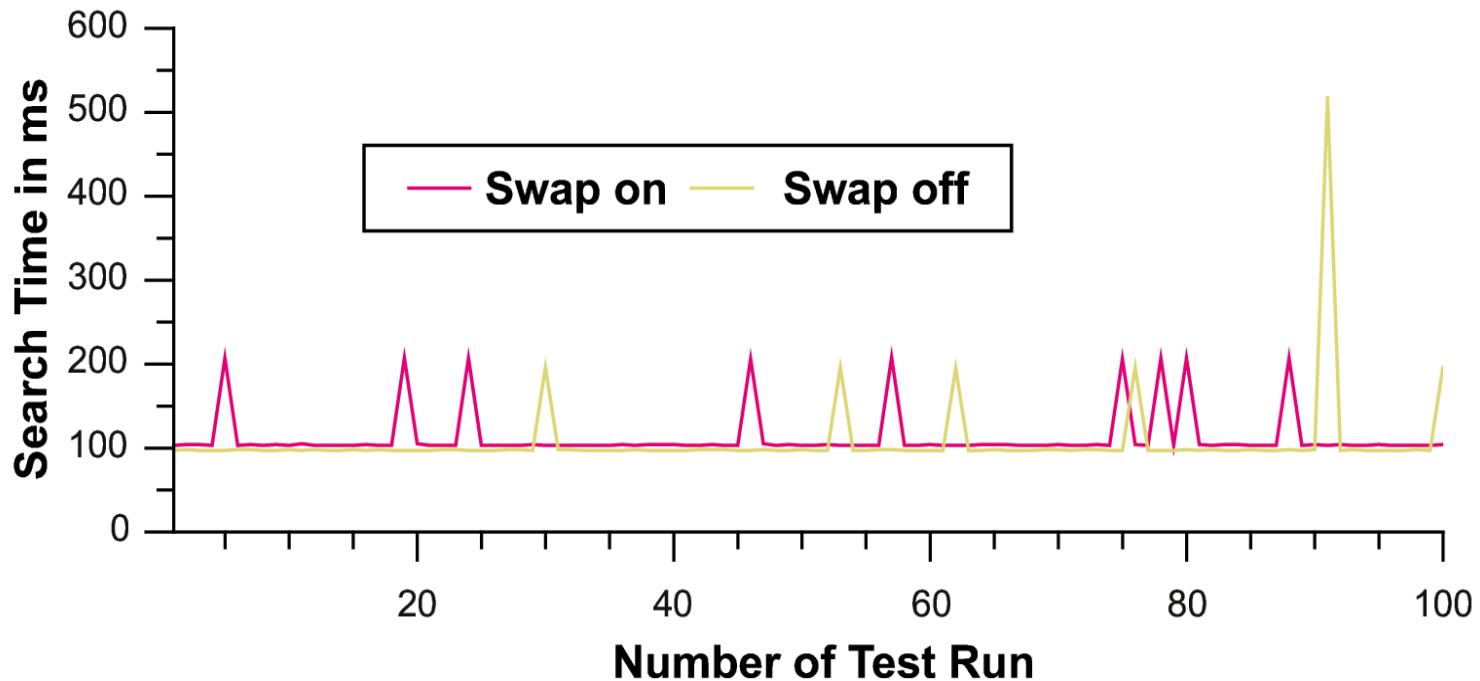
Evaluation – Several Keyboards, Windows Host



Evaluation – Several Keyboards, Linux Host



Evaluation – Swap File Behavior, Windows Host



Evaluation – ME Firmware Condition

- Different hooking strategies for Windows and Linux attacks
- Windows
 - *Local Manageability Service* driver
 - *AMT Status Tool*
 - *Manageability Developer Toolkit*
 - *Manageability Connector Tool*
- Linux
 - *Intel AMT Open-source Tools and Drivers*
 - *ME Status*
 - *ZTCLocalAgent*
- AMT webserver